BCH-hints

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1 Simulation von FEC über unterschiedliche Kanäle

1.1 Basic setup

```
In [1]: import numpy as np
from pprint import pprint as pp
%matplotlib notebook
from matplotlib import pyplot as plt
```

Die Variable behlist enthält eine Liste von Triples von (Paketlänge, Payload, maximal korrigierbare Fehler).

```
In [2]: # properties of bch as (n, k, t) triples
        bchlist = [(127, 127, 0)],
                 (127, 120, 1),
                 (127, 113, 2),
                 (127, 106, 3),
                 (127, 99, 4),
                 (127, 92, 5),
                 (127, 85, 6),
                 (127, 78, 7),
                 (127, 71, 9),
                 (127, 64, 10),
                 (127, 57, 11),
                    1
        packetlength = 127
        # not elegant: packetlength has to be identical for all BCH codes in the fo
        # let's also try longer packets:
        packetlength = 1023
        bchlist = [
            (1023, 1023, 0),
            (1023, 973, 5),
            (1023, 923, 10),
            (1023, 873, 15),
            (1023, 828, 20),
            (1023, 778, 25),
```

```
(1023, 728, 30),
(1023, 698, 35),
(1023, 648, 41),
(1023, 608, 45),
(1023, 573, 50),
]
# how many repetitions to run, for stochastic confidence?
# also: the more repetition, the more we can find rare events!
reps = 1000000
```

1.2 AWGN channel

Simple iteration over various SNR values

```
In [3]: snrdB = np.linspace(-5, 15, num=40)
    snr = 10**(snrdB/10.0)
    # bit error rate, assuming BPSK modulation.
    berlist = 0.5 * np.exp(-1*snr)
    # pp(snrdB)
    # pp(snr)
    # pp(berlist)
```

1.2.1 Paket error rate over AWGN

```
In [ ]: def num_errors_awgn(n, ber):
            """For a packet of length n and a given bit error rate,
            randomly generate number of errors in an AWGN channel"""
            return np.random.binomial(n, ber, 1)[0]
In [ ]: def compute_per(bchlist, berlist, reps=reps):
            """For every BCH code in bchlist, and every BER value in
            berlist, simpute reps many packet transmissions.
            A tranmission succeeds if it has fewer simulated errors than
            the considered BCH code can correct.
            .....
            per = \{\}
            for bch in bchlist:
                per[bch] = []
                for ber in berlist:
                    tmp = [num_errors_awgn(bch[0], ber) > bch[2] for i in range(rep
                    per[bch].append( (sum(tmp))/reps)
            return per
```

```
per = compute_per(bchlist, berlist)
In []: # Let's do a plot: BER as independent variables, for all BCH codes, show
        # packet error curves.
        # Convention: double-logarithmic plots over SNR and PER;
        # we do have logarithmic values for SNR already in the variable SNDdB,
        # so from a plotting tool perspective, this is only a semi-logarithmic plot
        plt.figure()
        for bch in bchlist:
            # Add plotting command here:
            ### BEGIN SOLUTION
            plt.semilogy(snrdB, per[bch])
            ### END SOLUTION
        plt.title("PER for different FEC schemes (AWGN)")
        plt.xlabel("SNR [dB]")
        plt.ylabel("PER")
        plt.show()
```

1.2.2 Throughput over AWGN

Let us also compute the effective throughput obtained for different BCH codes. On one hand, stronger code reduces error probability; on the other hand, it reduces payload length. So what is optimal code for a given SNR value?

```
In [ ]: throughput = { }
        for bch in bchlist:
            throughput[bch] = []
            for p in per[bch]:
                payload = bch[1]
                packetlength = bch[0]
                # Assign to variable tp the throughput obtained for
                # packet error rate p when there are payload many useful
                # bits in a paket of length packetlength (doh).
                # (It is ok to normalize this to a paketduration here since
                # all our codes have same length.
                # TODO: Think whether this needs extension!)
                ### BEGIN SOLUTION
                tp = payload/packetlength * (1-p)
                ### END SOLUTION
                throughput[bch].append(tp)
In [ ]: # plot throughput over SNR (logarithmic on SNR, natural unit on throughput,
        plt.figure()
        for bch in bchlist:
            plt.plot(snrdB, throughput[bch])
        plt.title("Throughput relative to uncoded transmission for different FEC so
```

```
plt.xlabel("SNR [dB]")
plt.ylabel("Througput")
plt.show()
```

1.3 FEC over a bursty channel

We switch from a simple AWGN scheme to a Gilbert-Elliot type bursty channel. First, the computation of number of errors in a frame is a bit more complicated. We define a class to hold state information: is the channel in a good or bad state?

```
In [ ]: class GilbertElliot():
            def __init__(self, berGood, berBad, gb, bg):
                self.ber = (berGood, berBad)
                self.M = ((1-qb, qb), (bq, 1-bq))
                self.state = 0
            def num_errors(self, n):
                """Simulate numer of errors in the next n bits."""
                errors = 0
                count = 0
                while count < n:</pre>
                    # how long do we stay in this state?
                    Pnextstate = self.M[self.state][1-self.state]
                    berstate = self.ber[self.state]
                    next_change = np.random.geometric(
                        Pnextstate, 1)[0]
                    # truncate to remaining packet length;
                    # only switch state if this falls indeed inside this packet
                    # Question: Explain why it would not be correct to change state
                    # ALWAYS, irrespective of this test!
                    if next_change <= n-count:</pre>
                        self.state = 1-self.state
                    else:
                        next_change = n-count
                    # how many errors within these many bits?
                    next_errors = np.random.binomial(
                        next_change, berstate)
                    errors += next_errors
                    count += next_change
                    # pp((Pnextstate, next_change, next_errors, count, errors))
```

```
return errors

def steady_state_ber(self):
    # steady states - why is this correct?
    gb = self.M[0][1]
    bg = self.M[1][0]
    tmp = gb + bg
    Pg = bg / tmp
    Pb = gb / tmp

    # pp((Pg, Pb))
    return Pg*self.ber[0] + Pb*self.ber[1]

def sim_steady_state_ber(self):
    n = 1000000
    err = self.num_errors(n)
    return err/n
```

Let's get an example channel object and look at its long-term steady-state, average bit error rate.

So it turns out that these parameters give us a pretty bad channel, but to illustrate the effects, that's ok)

1.3.1 A histogram for a GE channel

What is the distribution of number of errors in a packet under such a channel assumption?

```
In []: errors = [ge.num_errors(packetlength) for i in range(100000)]
    ## pp(errors)
    plt.figure()
    n, bins, patches= plt.hist(errors, bins=range(max(errors)))
    plt.show()
In []: pp(n)
    pp(bins)
    pp(sum(n))
    testber = sum([x*y for x, y in zip(n, bins)])/sum(n)/packetlength
    pp(testber)
```

Compare against an AWGN channel with the same BER:

```
In []: ge_ber = ge.steady_state_ber()
    errors = [num_errors_awgn(packetlength, ge_ber) for i in range(10000)]
    # pp(errors)
    plt.figure()
    n, bins, patches= plt.hist(errors, bins=range(max(errors)))
    plt.show()
In []: pp(n)
    pp(bins)
    testber = sum([x*y for x, y in zip(n, bins)])/sum(n)/packetlength
    pp(testber)
```

What is the interpretation here?

1.3.2 BEGIN SOLUTION

In an AWGN channel, there is a decent chance of seeing packets with no or very few errors. It is extremely unlikely to see very large numbers of errors.

In a bursty channel, on there other hand, there is a good chance to see even a lot of errors in a packet (and these packets will not be helped by FEC!). When comparing this to an AWGN channel of equivalent long-term BER, this leads to many packets that are error free. ### END SOLUTION

1.3.3 Use bursty channel with FEC codes

```
In [ ]: def compute_bursty_per(bchlist, ge):
            per_bursty = []
            per_awgn = []
            ss_ber = ge.steady_state_ber()
            for bch in bchlist:
                tmp = [ge.num_errors(bch[0]) > bch[2] for i in range(reps)]
                per_bursty.append(sum(tmp)/reps)
                tmp = [num_errors_awgn(bch[0], ss_ber) > bch[2] for i in range(reps
                per_awgn.append(sum(tmp)/reps)
            return (per_bursty, per_awqn)
        per2 = compute_bursty_per(bchlist, ge)
        pp(per2)
In []: # to plot, let's first get a good horizonatel axis:
        t = [bch[2] for bch in bchlist]
       plt.figure()
        plt.semilogy(t, per2[1], label="AWGN")
        plt.semilogy(t, per2[0], label="bursty")
```

```
plt.ylabel("PER after FEC")
       plt.title("PER for BCH AWGN and bursty channel")
       plt.xlabel("Number of correctable bits")
       plt.legend()
       plt.show()
In [ ]: # plot throughput
        throughput_awgn = [(1-per) * (bch[1]/bch[0])
                           for bch, per in zip(bchlist, per2[1])]
        throughput_bursty = [(1-per) * (bch[1]/bch[0])]
                             for bch, per in zip(bchlist, per2[0])]
       plt.figure()
       plt.plot(t, throughput_awgn, label="AWGN")
       plt.plot(t, throughput_bursty, label="bursty")
        plt.title("Throughput for BCH AWGN and bursty channel")
       plt.ylabel("Throughput after FEC")
       plt.xlabel("Number of correctable bits")
       plt.legend()
       plt.show()
```

2 Delay

And what about delay characteristics? We save this for another time